ITT8060: Advanced Programming (in F#) Lecture 5: Records. Sets and maps as abstract data types

Juhan Ernits, Hendrik Maarand and Ian Erik Varatalu based on slides by Michael R. Hansen

Department of Software Science, Tallinn University of Technology

04/10/2023

(日)

Overview

Sets and Maps as abstract data types

- Useful in the modelling and solution of many problems
- Many similarities with the list library

Recommendation: Use these libraries whenever it is appropriate.

(ロ)、

The set concept (1)

A set (in mathematics) is a collection of element like

```
\{Bob, Bill, Ben\}, \{1, 3, 5, 7, 9\}, \mathbb{N}, and \mathbb{R}
```

the sequence in which elements are enumerated is of no concern, and
 repetitions among members of a set is of no concern either

It is possible to decide whether a given value is in the set.

Alice \notin {Bob, Bill, Ben} and $7 \in$ {1, 3, 5, 7, 9}

< □ > < 同 > < 三 > < 三 > < 三 > < ○ < ○ </p>

The empty set containing no element is written $\{\}$ or \emptyset .

The set concept (1)

A set (in mathematics) is a collection of element like

```
\{Bob, Bill, Ben\}, \{1, 3, 5, 7, 9\}, \mathbb{N}, and \mathbb{R}
```

the sequence in which elements are enumerated is of no concern, and
 repetitions among members of a set is of no concern either

It is possible to decide whether a given value is in the set.

Alice \notin {Bob, Bill, Ben} and $7 \in$ {1, 3, 5, 7, 9}

< □ > < 同 > < 三 > < 三 > < 三 > < ○ < ○ </p>

The empty set containing no element is written $\{\}$ or \emptyset .

The set concept (1)

A set (in mathematics) is a collection of element like

```
\{Bob, Bill, Ben\}, \{1, 3, 5, 7, 9\}, \mathbb{N}, and \mathbb{R}
```

the sequence in which elements are enumerated is of no concern, and
 repetitions among members of a set is of no concern either

It is possible to decide whether a given value is in the set.

Alice \notin {Bob, Bill, Ben} and $7 \in$ {1, 3, 5, 7, 9}

< □ > < 同 > < 三 > < 三 > < 三 > < ○ < ○ </p>

The empty set containing no element is written $\{\}$ or \emptyset .

The sets concept (2)

A set A is a *subset* of a set B, written $A \subseteq B$, if all the elements of A are also elements of B, for example

 $\{Ben, Bob\} \subseteq \{Bob, Bill, Ben\}$ and $\{1, 3, 5, 7, 9\} \subseteq \mathbb{N}$

Two sets A and B are equal, if they are both subsets of each other:

A = B if and only if $A \subseteq B$ and $B \subseteq A$

i.e. two sets are equal if they contain exactly the same elements.

The subset of a set A which consists of those elements satisfying a predicate p can be expressed using a *set-comprehension*:

 $\{x \in A \mid p(x)\}$

For example:

 $\{1, 3, 5, 7, 9\} = \{x \in \mathbb{N} \mid \text{odd}(x) \text{ and } x < 11\}$

・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・

The sets concept (2)

A set A is a *subset* of a set B, written $A \subseteq B$, if all the elements of A are also elements of B, for example

 $\{Ben, Bob\} \subseteq \{Bob, Bill, Ben\}$ and $\{1, 3, 5, 7, 9\} \subseteq \mathbb{N}$

Two sets A and B are equal, if they are both subsets of each other:

A = B if and only if $A \subseteq B$ and $B \subseteq A$

i.e. two sets are equal if they contain exactly the same elements.

The subset of a set A which consists of those elements satisfying a predicate p can be expressed using a *set-comprehension*:

 $\{x \in A \mid p(x)\}$

For example:

 $\{1, 3, 5, 7, 9\} = \{x \in \mathbb{N} \mid \text{odd}(x) \text{ and } x < 11\}$

・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・

The sets concept (2)

A set A is a *subset* of a set B, written $A \subseteq B$, if all the elements of A are also elements of B, for example

 $\{Ben, Bob\} \subseteq \{Bob, Bill, Ben\}$ and $\{1, 3, 5, 7, 9\} \subseteq \mathbb{N}$

Two sets A and B are equal, if they are both subsets of each other:

A = B if and only if $A \subseteq B$ and $B \subseteq A$

i.e. two sets are equal if they contain exactly the same elements.

The subset of a set *A* which consists of those elements satisfying a predicate *p* can be expressed using a *set-comprehension*:

 $\{x \in A \mid p(x)\}$

For example:

 $\{1,3,5,7,9\} = \{x \in \mathbb{N} \mid \text{odd}(x) \text{ and } x < 11\}$

The set concept (3)

Some standard operations on sets:



Figure: Venn diagrams for (a) union, (b) intersection and (c) difference

For example

 $\{Bob, Bill, Ben\} \cup \{Alice, Bill, Ann\} = \{Alice, Ann, Bob, Bill, Ben\} \\ \{Bob, Bill, Ben\} \cap \{Alice, Bill, Ann\} = \{Bill\} \\ \{Bob, Bill, Ben\} \setminus \{Alice, Bill, Ann\} = \{Bob, Ben\}$

< □ > < 同 > < 三 > < 三 > < 三 > < ○ < ○ </p>

The set concept (3)

Some standard operations on sets:



Figure: Venn diagrams for (a) union, (b) intersection and (c) difference

For example

$$\begin{split} & \{ \text{Bob}, \text{Bill}, \text{Ben} \} \cup \{ \text{Alice}, \text{Bill}, \text{Ann} \} &= \{ \text{Alice}, \text{Ann}, \text{Bob}, \text{Bill}, \text{Ben} \} \\ & \{ \text{Bob}, \text{Bill}, \text{Ben} \} \cap \{ \text{Alice}, \text{Bill}, \text{Ann} \} &= \{ \text{Bill} \} \\ & \{ \text{Bob}, \text{Bill}, \text{Ben} \} \setminus \{ \text{Alice}, \text{Bill}, \text{Ann} \} &= \{ \text{Bob}, \text{Ben} \} \end{split}$$

An abstract Data Type: A type together with a collection of operations, where

◆□▶ ◆□▶ ▲□▶ ▲□▶ □ のQ@

the representation of values is hidden.

An abstract data type for sets must have:

- Operations to generate sets from the elements. Why?
- Operations to extract the elements of a set. Why?
- Standard operations on sets.

An abstract Data Type: A type together with a collection of operations, where

(ロ) (同) (三) (三) (三) (○) (○)

the representation of values is hidden.

An abstract data type for sets must have:

- Operations to generate sets from the elements. Why?
- Operations to extract the elements of a set. Why?
- Standard operations on sets.

Sets in F#

The Set library of F# supports finite sets. An efficient implementation is based on a balanced binary tree.

Examples:

```
set ["Bob"; "Bill"; "Ben"];;
val it : Set<string> = set ["Ben"; "Bill"; "Bob"]
set [3; 1; 9; 5; 7; 9; 1];;
val it : Set<int> = set [1; 3; 5; 7; 9]
```

Equality of two sets is tested in the usual manner:

```
set["Bob";"Bill";"Ben"] = set["Bill";"Ben";"Bill";"Bob"];;
val it : bool = true
```

Sets are order on the basis of a lexicographical ordering:

compare (set ["Ann";"Jane"]) (set ["Bill";"Ben";"Bob"]);;
val it : int = -1

▲□▶▲□▶▲□▶▲□▶ □ のQ@

Selected operations (1)

Notice that minElement is well-defined due to the ordering:

Set.minElement (Set.ofList ["Bob"; "Bill"; "Ben"]);; val it : string = "Ben"

▲□▶ ▲□▶ ▲□▶ ▲□▶ = 三 のへで

Selected operations (1)

Notice that minElement is well-defined due to the ordering:

Set.minElement (Set.ofList ["Bob"; "Bill"; "Ben"]);;
val it : string = "Ben"

Selected operations (2)

- vunion: Set<'a> -> Set<'a> -> Set<'a>,
 where union A B = A ∪ B
- intersect: Set<'a> -> Set<'a> -> Set<'a>,
 where intersect A B = A ∩ B
- difference: Set<'a> -> Set<'a> -> Set<'a>, where difference A B = A \ B
- ▶ exists: ('a -> bool) -> Set<'a> -> bool, where exists $p A = \exists x \in A.p(x)$
- ▶ forall: ('a -> bool) -> Set<'a> -> bool, where forall $p A = \forall x \in A.p(x)$
- fold: ('a -> 'b -> 'a) -> 'a -> Set<'b> -> 'a,
 where

fold
$$f a \{b_0, b_1, \dots, b_{n-2}, b_{n-1}\}$$

= $f(f(f(\dots f(f(a, b_0), b_1), \dots), b_{n-2}), b_{n-1})$

These work similar to their List siblings, e.g.

Set.fold (-) 0 (set [1; 2; 3]) = ((0-1)-2)-3 = -6

where the ordering is exploited.

Selected operations (2)

- union: Set<'a> -> Set<'a> -> Set<'a>,
 where union A B = A ∪ B
- intersect: Set<'a> -> Set<'a> -> Set<'a>,
 where intersect A B = A ∩ B
- difference: Set<'a> -> Set<'a> -> Set<'a>, where difference A B = A \ B
- ▶ exists: ('a -> bool) -> Set<'a> -> bool, where exists $p A = \exists x \in A.p(x)$
- ▶ forall: ('a -> bool) -> Set<'a> -> bool, where forall $p A = \forall x \in A.p(x)$
- fold: ('a -> 'b -> 'a) -> 'a -> Set<'b> -> 'a,
 where

fold
$$f a \{b_0, b_1, \dots, b_{n-2}, b_{n-1}\}$$

= $f(f(f(\cdots f(f(a, b_0), b_1), \dots), b_{n-2}), b_{n-1})$

These work similar to their List siblings, e.g.

Set.fold (-) 0 (set [1; 2; 3]) = ((0-1)-2)-3 = -6

where the ordering is exploited.

Example: Map Coloring (1)

Maps and colors are modelled in a more natural way using sets:

```
type country = string;;
type map = Set<country*country>;;
type color = Set<country>;;
type coloring = Set<color>;;
```

WHY?

◆□▶ ◆□▶ ◆□▶ ◆□▶ ● ● ● ●

```
Iwo countries c<sub>1</sub>, c<sub>2</sub> are neighbors in a map m,
if either (c<sub>1</sub>, c<sub>2</sub>) \in m or (c<sub>2</sub>, c<sub>1</sub>) \in m:
let areNb c1 c2 m =
   Set.contains (c1, c2) m || Set.contains (c2, c1) m;;
```

Color *col* and be extended by a country *c* given map *m*, if for every country *c'* in *col*: *c* and *c'* are not neighbours in *m*

```
let canBeExtBy m col c =
   Set.forall (fun c' -> not (areNb c' c m)) col;;
```

Example: Map Coloring (1)

Maps and colors are modelled in a more natural way using sets:

```
type country = string;;
type map = Set<country*country>;;
type color = Set<country>;;
type coloring = Set<color>;;
```

```
WHY?
```

◆□▶ ◆□▶ ▲□▶ ▲□▶ □ のQ@

Two countries c_1, c_2 are neighbors in a map m, if either $(c_1, c_2) \in m$ or $(c_2, c_1) \in m$: let areNb c1 c2 m =

Set.contains (c1,c2) m || Set.contains (c2,c1) m;;

Color *col* and be extended by a country *c* given map *m*, if for every country *c'* in *col*: *c* and *c'* are not neighbours in *m*

```
let canBeExtBy m col c =
   Set.forall (fun c' -> not (areNb c' c m)) col;;
```

Example: Map Coloring (1)

Maps and colors are modelled in a more natural way using sets:

```
type country = string;;
type map = Set<country*country>;;
type color = Set<country>;;
type coloring = Set<color>;;
WHY?
Two countries c_1, c_2 are neighbors in a map m,
if either (c_1, c_2) \in m or (c_2, c_1) \in m:
let areNb c1 c2 m =
Set.contains (c1, c2) m || Set.contains (c2, c1) m;;
```

Color *col* and be extended by a country *c* given map *m*, if for every country c' in *col*: *c* and c' are not neighbours in *m*

```
let canBeExtBy m col c =
   Set.forall (fun c' -> not (areNb c' c m)) col;;
```

(ロ) (同) (三) (三) (三) (○) (○)

Example: Map Coloring (2)

The function

```
extColoring: map -> coloring -> country -> coloring
```

is declared as a recursive function over the coloring:

WHY not use a fold function?

```
let rec extColoring m cols c =
    if Set.isEmpty cols
    then Set.singleton (Set.singleton c)
    else let col = Set.minElement cols
        let cols' = Set.remove col cols
        if canBeExtBy m col c
        then Set.add (Set.add c col) cols'
        else Set.add col (extColoring m cols' c);;
```

Notice similarity to a list recursion:

- base case [] corresponds to the empty set
- for a recursive case x::xs, the head x corresponds to the minimal element col and the tail xs corresponds to the "rests" set cols'

The list-based version is more efficient (why?) and more readable.

Example: Map Coloring (2)

The function

```
extColoring: map -> coloring -> country -> coloring
```

is declared as a recursive function over the coloring:

```
WHY not use a fold function?
```

```
let rec extColoring m cols c =
    if Set.isEmpty cols
    then Set.singleton (Set.singleton c)
    else let col = Set.minElement cols
        let cols' = Set.remove col cols
        if canBeExtBy m col c
        then Set.add (Set.add c col) cols'
        else Set.add col (extColoring m cols' c);;
```

Notice similarity to a list recursion:

- base case [] corresponds to the empty set
- for a recursive case x::xs, the head x corresponds to the minimal element col and the tail xs corresponds to the "rests" set cols'

The list-based version is more efficient (why?) and more readable.

Example: Map Coloring (2)

The function

```
extColoring: map -> coloring -> country -> coloring
```

is declared as a recursive function over the coloring:

```
WHY not use a fold function?
```

```
let rec extColoring m cols c =
    if Set.isEmpty cols
    then Set.singleton (Set.singleton c)
    else let col = Set.minElement cols
        let cols' = Set.remove col cols
        if canBeExtBy m col c
        then Set.add (Set.add c col) cols'
        else Set.add col (extColoring m cols' c);;
```

Notice similarity to a list recursion:

- base case [] corresponds to the empty set
- for a recursive case x::xs, the head x corresponds to the minimal element col and the tail xs corresponds to the "rests" set cols'

The list-based version is more efficient (why?) and more readable.

Example: Map Coloring (3)

A set of countries is obtained from a map by the function:

```
countries: map -> Set<country>
```

that is based on repeated insertion of the countries into a set:

```
let countries m =
   Set.fold
   (fun set (c1,c2) -> Set.add c1 (Set.add c2 set))
    Set.empty
   m;;
```

The function

```
colCntrs: map -> Set<country> -> coloring
```

is based on repeated insertion of countries in colorings using the extColoring function:

```
let colCntrs m cs = Set.fold (extColoring m) Set.empty cs;
```

▲□▶▲□▶▲□▶▲□▶ □ のQ@

Example: Map Coloring (3)

A set of countries is obtained from a map by the function:

```
countries: map -> Set<country>
```

that is based on repeated insertion of the countries into a set:

```
let countries m =
   Set.fold
   (fun set (c1,c2) -> Set.add c1 (Set.add c2 set))
    Set.empty
   m;;
```

The function

```
colCntrs: map -> Set<country> -> coloring
```

is based on repeated insertion of countries in colorings using the extColoring function:

```
let colCntrs m cs = Set.fold (extColoring m) Set.empty cs;;
```

▲□▶▲□▶▲□▶▲□▶ □ のQ@

The function that creates a coloring from a map is declared using functional composition:

The map concept

A map from a set A to a set B is a finite subset A' of A together with a function m defined on $A': m: A' \rightarrow B$.

The set A' is called the *domain* of m: dom m = A'.

A map *m* can be described in a tabular form:



- An element a_i in the set A' is called a key
- A pair (a_i, b_i) is called an *entry*, and
- \blacktriangleright *b_i* is called the *value* for the key *a_i*.

We denote the sets of entries of a map as follows:

entriesOf(m) = {(a_0, b_0), ..., (a_{n-1}, b_{n-1})}

(日) (日) (日) (日) (日) (日) (日)

Selected map operations in F#

- ofList: ('a*'b) list -> Map<'a,'b>
 ofList [(a₀, b₀);...; (a_{n-1}, b_{n-1})] = m
- add: 'a -> 'b -> Map<'a, 'b> -> Map<'a, 'b> add a b m = m', where m' is obtained m by overriding m with the entry (a, b)

- ▶ tryFind: 'a -> Map<'a, 'b> -> 'b option tryFind a m = Some (m(a)), if a ∈ dom m; None otherwise
- foldBack: ('a->'b->'c->'c) -> Map<'a,'b> -> 'c -> 'c
 foldBack f m c = f a₀ b₀ (f a₁ b₁ (f ... (f a_{n-1} b_{n-1} c)···))

A D F A 同 F A E F A E F A Q A

A few examples

An entry can be added to a map using add and the value for a key in a map is retrieved using either find or tryFind:

```
Map.find "a2" reg1;;
val it : string * int = ("herring", 4)
```

```
Map.tryFind "a2" reg1;;
val it : (string * int) option = Some ("herring", 4)
```

We can extract the list of article codes and prices for a given register using the fold functions for maps:

A D F A 同 F A E F A E F A Q A

This and other higher-order functions are similar to their List and Set siblings.

Example: Cash register (1)

```
type articleCode = string;;
type articleName = string;;
type noPieces = int;;
type price = int;;
type info = noPieces * articleName * price;;
type infoseq = info list;;
type bill = infoseq * price;;
```

The natural model of a register is using a map:

```
type register = Map<articleCode, articleName*price>;;
since an article code is a unique identification of an article.
First version:
    type item = noPieces * articleCode;;
```

```
pe purchase     = item list;
```

Example: Cash register (1)

```
type articleCode = string;;
type articleName = string;;
type noPieces = int;;
type price = int;;
type info = noPieces * articleName * price;;
type infoseq = info list;;
type bill = infoseq * price;;
```

The natural model of a register is using a map:

```
type register = Map<articleCode, articleName*price>;;
```

since an article code is a unique identification of an article.

First version:

type item = noPieces * articleCode;;
type purchase = item list;;

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ ─臣 ─のへで

Example: Cash register (1) - a recursive program

```
exception FindArticle;;
(* makebill: register -> purchase -> bill *)
let rec makeBill reg = function
    | [] -> ([],0)
   | (np,ac)::pur ->
       match Map.tryFind ac req with
        | None
                        -> raise FindArticle
        | Some(aname,aprice) ->
           let tprice = np*aprice
           let (infos, sumbill) = makeBill reg pur
            ((np,aname,tprice)::infos, tprice+sumbill);;
let pur = [(3,"a2"); (1,"a1")];;
makeBill reg1 pur;;
val it : (int * string * int) list * int =
  ([(3, "herring", 12); (1, "cheese", 25)], 37)
```

< □ > < 同 > < 三 > < 三 > < 三 > < ○ < ○ </p>

the lookup in the register is managed by a Map.tryFind

Example: Cash register (1) - a recursive program

```
exception FindArticle;;
(* makebill: register -> purchase -> bill *)
let rec makeBill reg = function
   | [] -> ([],0)
   | (np,ac)::pur ->
       match Map.tryFind ac req with
        | None
                        -> raise FindArticle
        | Some(aname,aprice) ->
           let tprice = np*aprice
           let (infos, sumbill) = makeBill reg pur
            ((np,aname,tprice)::infos, tprice+sumbill);;
let pur = [(3,"a2"); (1,"a1")];;
makeBill reg1 pur;;
val it : (int * string * int) list * int =
  ([(3, "herring", 12); (1, "cheese", 25)], 37)
```

the lookup in the register is managed by a Map.tryFind

Example: Cash register (2) - using List.foldBack

```
let makeBill' reg pur =
    let f (np,ac) (infos,billprice)
        = let (aname, aprice) = Map.find ac reg
        let tprice = np*aprice
        ((np,aname,tprice)::infos, tprice+billprice)
    List.foldBack f pur ([],0);;
makeBill' reg1 pur;;
```

▲□▶▲□▶▲□▶▲□▶ □ のQ@

```
val it : (int * string * int) list * int =
  ([(3, "herring", 12); (1, "cheese", 25)], 37)
```

the recursion is handled by List.foldBack

the exception is handled by Map.find

Example: Cash register (2) - using List.foldBack

```
let makeBill' reg pur =
    let f (np,ac) (infos,billprice)
        = let (aname, aprice) = Map.find ac reg
        let tprice = np*aprice
        ((np,aname,tprice)::infos, tprice+billprice)
    List.foldBack f pur ([],0);;
makeBill' reg1 pur;;
val it : (int * string * int) list * int =
```

◆□▶ ◆□▶ ▲□▶ ▲□▶ □ のQ@

```
([(3, "herring", 12); (1, "cheese", 25)], 37)
```

the recursion is handled by List.foldBack

the exception is handled by Map.find

Example: Cash register (2) - using List.foldBack

```
let makeBill' reg pur =
    let f (np,ac) (infos,billprice)
        = let (aname, aprice) = Map.find ac reg
        let tprice = np*aprice
        ((np,aname,tprice)::infos, tprice+billprice)
    List.foldBack f pur ([],0);;
makeBill' regl pur;;
val it : (int * string * int) list * int =
    ([(3, "herring", 12); (1, "cheese", 25)], 37)
```

◆□▶ ◆□▶ ▲□▶ ▲□▶ □ のQ@

- the recursion is handled by List.foldBack
- the exception is handled by Map.find

Example: Cash register (2) - using maps for purchases

The purchase: 3 herrings, one piece of cheese, and 2 herrings, is the same as a purchase of one piece of cheese and 5 herrings.

A purchase associated number of pieces with article codes:

```
type purchase = Map<articleCode,noPieces>;;
```

A bill is produced by folding a function over a map-purchase:

```
let makeBill'' reg pur =
   let f ac np (infos,billprice)
        = let (aname, aprice) = Map.find ac reg
        let tprice = np*aprice
        ((np,aname,tprice)::infos, tprice+billprice)
   Map.foldBack f pur ([],0);;
```

```
let purMap = Map.ofList [("a2",3); ("a1",1)];;
val purMap : Map<string,int> = map [("a1", 1); ("a2", 3)]
```

```
makeBill'' reg1 purMap;;
val it = ([(1, "cheese", 25); (3, "herring", 12)], 37)
```

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ - 三 - のへぐ

Example: Cash register (2) - using maps for purchases

The purchase: 3 herrings, one piece of cheese, and 2 herrings, is the same as a purchase of one piece of cheese and 5 herrings.

A purchase associated number of pieces with article codes:

type purchase = Map<articleCode,noPieces>;;

A bill is produced by folding a function over a map-purchase:

```
let makeBill'' reg pur =
   let f ac np (infos,billprice)
        = let (aname, aprice) = Map.find ac reg
        let tprice = np*aprice
        ((np,aname,tprice)::infos, tprice+billprice)
   Map.foldBack f pur ([],0);;
```

```
let purMap = Map.ofList [("a2",3); ("a1",1)];;
val purMap : Map<string,int> = map [("a1", 1); ("a2", 3)]
```

```
makeBill'' reg1 purMap;;
val it = ([(1, "cheese", 25); (3, "herring", 12)], 37)
```

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ - 三 - のへぐ

Example: Cash register (2) - using maps for purchases

The purchase: 3 herrings, one piece of cheese, and 2 herrings, is the same as a purchase of one piece of cheese and 5 herrings.

A purchase associated number of pieces with article codes:

```
type purchase = Map<articleCode,noPieces>;;
```

A bill is produced by folding a function over a map-purchase:

```
let makeBill'' reg pur =
    let f ac np (infos,billprice)
        = let (aname, aprice) = Map.find ac reg
        let tprice = np*aprice
        ((np,aname,tprice)::infos, tprice+billprice)
    Map.foldBack f pur ([],0);;
let purMap = Map.ofList [("a2",3); ("a1",1)];;
val purMap : Map<string,int> = map [("a1", 1); ("a2", 3)]
makeBill'' reg1 purMap;;
val it = ([(1, "cheese", 25); (3, "herring", 12)], 37)
```

Summary

- The concepts of sets and maps.
- Fundamental operations on sets and maps.

(ロ)、

Applications of sets and maps.